

Heap Sort

Heap Sort is a **comparison-based sorting algorithm** that uses a special data structure called a **Heap**.

It is based on the idea of using a **Binary Heap** (usually a Max Heap) to repeatedly extract the maximum element and build a sorted array.

Heap Sort is efficient and works in-place, but unlike Merge Sort and Quick Sort, it is not stable.

- A Heap is a complete binary tree that satisfies the Heap Property.
- Max Heap: The parent node is always greater than or equal to its children.
- Min Heap: The parent node is always smaller than or equal to its children.

Heap Sort generally uses a **Max Heap** to sort numbers in ascending order.

Steps in Heap Sort

- 1. **Build Max Heap** → Convert the array into a Max Heap (largest element at the root).
- Swap Root with Last Element → Swap the maximum element (root) with the last element of the heap.
- 3. **Reduce Heap Size & Heapify** → Exclude the last (sorted) element and heapify the remaining elements to maintain heap property.
- Repeat → Continue until all elements are sorted.

Example

Sort the array: [4, 10, 3, 5, 1]

• Step 1: Build Max Heap → [10, 5, 3, 4, 1]

- Step 2: Swap root with last → [1, 5, 3, 4, 10]
- Step 3: Heapify → [5, 4, 3, 1, 10]
- Step 4: Repeat process → [1, 3, 4, 5, 10]

Final Sorted Array: [1, 3, 4, 5, 10]

Time Complexity

- Building Heap: 0(n)
- Heapify (per element): $0(\log n)$
- Overall: 0(n log n)
- Best Case: 0(n log n)
- Average Case: 0(n log n)
- Worst Case: 0(n log n)

Unlike Quick Sort, Heap Sort's worst case is not bad.

Space Complexity

0(1) → In-place sorting (only requires constant extra memory).

Advantages of Heap Sort

- Guaranteed O(n log n) time complexity.
- In-place sorting (no extra space like Merge Sort).





• Useful for Priority Queue operations (using heaps).

Disadvantages of Heap Sort

- Not Stable: Equal elements may not maintain their relative order.
- In practice, Heap Sort is often slower than Quick Sort due to cache inefficiency.

